

Optimal Student Sectioning at Niederrhein University of Applied Sciences

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Niederrhein University of Applied Sciences - Institute for Pattern Recognition,
Faculty of Electrical Engineering and Computer Science

OR 2019

Time table with parallel group sessions

	8:00 9:00	9:00 10:00	10:00 11:00	11:00 12:00	12:00 13:00	13:00 14:00	14:00 15:00	15:00 16:00	16:00 17:00	17:00 18:00	18:00 19:00	19:00 20:00	20:00 21:00
Monday			ENG F Dj		BSY V Po	MAR V Gp	MAR U Gp						
	OOA P St						OOA T ShI						
	OOA P Dv												
	BSY P Ni												
Tuesday	OOA P Gb			ALD V Ue	ALD U Ue		MAR U Gp						
	BSY P Ne				TEI2 U Na		BSY U Po						
	BSY P Po				BSY U Po		MINT T ShE						
	TEI2 P Na												
	OOA P Qi												
Wednesday	OOA U St	MA2 V Tp	OOA V St		OOA U St		ALD U Ue						
	MA2 U Tp		OOA V Dv		ALD U Ue		OOA U Dv						
	ENG F Nn												
Thursday	TEI2 V Ha	OOA V St		MA2 T ShI2	ENG F Ge		OOA U Dv		MINT T ShI				
		OOA V Dv			MA2 U Su		MA2 U Su						
					ALD U Ue								
					ENG F Db								
Friday	MA2 V Tp	MA2 U Tp		TEI2 V Ha	TEI2 P Ha								
		TEI2 U Ha			OOA P Le								

Model and Notation

student $s \in [S] := \{1, 2, \dots, S\}$ is enrolled for a module $k \in [N]$,
i.e. $c_{s,k} = 1$

$$b_{k,j,s} = 1$$

group	$G_{k,1}$	$G_{k,2}$	$G_{k,3}$...	$G_{k,j}$...	G_{k,n_k}
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$$a_{k,j,i,l} = 1$$

time slot	$T_{k,1}$	$T_{k,2}$...	$T_{k,i}$...	T_{k,m_k}
time sub-slot	$T_{k,1,1}$	$T_{k,2,1}$...	$T_{k,i,1}$...	$T_{k,m_k,1}$
	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots
	$T_{k,1,l}$	$T_{k,2,l}$...	$T_{k,i,l}$...	$T_{k,m_k,l}$
	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots
	$T_{k,1,p_k}$	$T_{k,2,p_k}$...	T_{k,i,p_k}	...	T_{k,m_k,p_k}

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time sub-slot	$T_{k,1,1}$	$T_{k,2,1}$...	$T_{k,i,1}$...	$T_{k,m_k,1}$
	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots
	$T_{k,1,l}$	$T_{k,2,l}$...	$T_{k,i,l}$...	$T_{k,m_k,l}$
	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots
	$T_{k,1,p_k}$	$T_{k,2,p_k}$...	T_{k,i,p_k}	...	T_{k,m_k,p_k}

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time sub-slot	$T_{k,1,1}$	$T_{k,2,1}$...	$T_{k,i,1}$...	$T_{k,m_k,1}$
	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots
	$T_{k,1,l}$	$T_{k,2,l}$...	$T_{k,i,l}$...	$T_{k,m_k,l}$
	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots
	$T_{k,1,p_k}$	$T_{k,2,p_k}$...	T_{k,i,p_k}	...	T_{k,m_k,p_k}

Restrictions (1)

For each group $G_{k,j}$, binary variables $a_{k,j,i,l}$ are used to assign a time sub-slot:

$$\sum_{i=1}^{m_k} \sum_{l=1}^{p_k} a_{k,j,i,l} = 1. \quad (1)$$

Every sub-slot $T_{k,i,l}$ has to be assigned at most once to a group ($k \in [N]$, $i \in [m_k]$, $l \in [p_k]$):

$$\sum_{j=1}^{n_k} a_{k,j,i,l} \leq 1. \quad (2)$$

Only $q_{k,i}$ groups may be scheduled for a time slot $T_{k,i}$, i.e. ($k \in [N]$, $i \in [m_k]$)

$$\sum_{j=1}^{n_k} \sum_{l=1}^{p_k} a_{k,j,i,l} \leq q_{k,i}. \quad (3)$$

Restrictions (2)

We use a matrix $C \in \{0, 1\}^{S \times N}$ to describe whether a student has enrolled for a module. Thereby, $c_{s,k} = 1$ indicates that student s has chosen module k .

Student s is in group j of module k iff $b_{k,j,s} = 1$.

For $k \in [N]$ and $s \in [S]$, we get condition

$$\sum_{j=1}^{n_k} b_{k,j,s} = c_{s,k}. \quad (4)$$

Time sub-slots

A group of module k assigned to a sub time-slot $T_{k,i,l}$ of time slot $T_{k,i}$ is taught depending on frequency p_k in following weeks:

frequency p_k	week 1	week 2	week 3	week 4
1	$T_{k,i,1}$	$T_{k,i,1}$	$T_{k,i,1}$	$T_{k,i,1}$
2	$T_{k,i,1}$	$T_{k,i,2}$	$T_{k,i,1}$	$T_{k,i,2}$
4	$T_{k,i,1}$	$T_{k,i,2}$	$T_{k,i,3}$	$T_{k,i,4}$

Time sub-slots

A group of module k assigned to a sub time-slot $T_{k,i,l}$ of time slot $T_{k,i}$ is taught depending on frequency p_k in following weeks:

frequency p_k	week 1	week 2	week 3	week 4
1	$T_{k,i,1}$	$T_{k,i,1}$	$T_{k,i,1}$	$T_{k,i,1}$
2	$T_{k,i,1}$	$T_{k,i,2}$	$T_{k,i,1}$	$T_{k,i,2}$
4	$T_{k,i,1}$	$T_{k,i,2}$	$T_{k,i,3}$	$T_{k,i,4}$

A student can be member of groups of different modules that are taught in overlapping time slots but in different weeks. For example let $p_{k_1} = 2$, $p_{k_2} = 4$, and $p_{k_4} = 4$:

week 1	week 2	week 3	week 4
$T_{k_1,i_1,1}$		$T_{k_1,i_1,1}$	
	$T_{k_2,i_2,2}$		
			$T_{k_3,i_3,4}$

Time sub-slots with conflicts

Conflicting sub-slots are calculated in advance:

- | | | | |
|-------------------|---|---|---|
| 1 | 2 | 3 | 4 |
| $T_{k_1, i_1, 1}$ | X | X | X |

Let a student be assigned to a group that attends a weekly sub-slot $T_{k_1, i_1, 1}$ of module k_1 .

⇒ He cannot attend other time-overlapping sub-slots.

Time sub-slots with conflicts

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|-------------------|---|---|---|
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⇒ He cannot attend other time-overlapping sub-slots.

- | 1 | 2 | 3 | 4 |
|-------------------|-------------------|---|---|
| $T_{k_1, i_1, 1}$ | $T_{k_1, i_1, 2}$ | X | X |

A student attends a bi-weekly sub-slot $T_{k_1, i_1, l}$.

⇒ He cannot attend an overlapping sub-slot $T_{k_2, i_2, l}$ if it is bi-weekly or four-weekly, and $T_{k_2, i_2, l+2}$ if it is four-weekly.

Time sub-slots with conflicts

Conflicting sub-slots are calculated in advance:

- | 1 | 2 | 3 | 4 |
|-------------------|---|---|---|
| $T_{k_1, i_1, 1}$ | X | X | X |

Let a student be assigned to a group that attends a weekly sub-slot $T_{k_1, i_1, 1}$ of module k_1 .

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- | 1 | 2 | 3 | 4 |
|-------------------|-------------------|---|---|
| $T_{k_1, i_1, 1}$ | $T_{k_1, i_1, 2}$ | X | X |

A student attends a bi-weekly sub-slot $T_{k_1, i_1, l}$.

⇒ He cannot attend an overlapping sub-slot $T_{k_2, i_2, l}$ if it is bi-weekly or four-weekly, and $T_{k_2, i_2, l+2}$ if it is four-weekly.

- | 1 | 2 | 3 | 4 |
|-------------------|-------------------|-------------------|-------------------|
| $T_{k_1, i_1, 1}$ | $T_{k_1, i_1, 2}$ | $T_{k_1, i_1, 3}$ | $T_{k_1, i_1, 4}$ |

A student is taught in a four-weekly time sub-slot $T_{k_1, i_1, l}$.

⇒ He must not be assigned to overlapping four-weekly sub-slots $T_{k_2, i_2, l}$.

Restrictions (3)

For all colliding pairs $(T_{k_1, i_1, l_1}, T_{k_2, i_2, l_2})$ of time sub-slots, all groups $j_1 \in [n_{k_1}]$, $j_2 \in [n_{k_2}]$ and all students $s \in [S]$,

$$a_{k_1, j_1, i_1, l_1} + b_{k_1, j_1, s} + a_{k_2, j_2, i_2, l_2} + b_{k_2, j_2, s} \leq 3 \quad (5)$$

has to be fulfilled.

Further restrictions are used to model rules for part-time and dual education students, car pools, learning groups, etc.

Objective function: groups of nearly equal size

Difference of sizes of groups j_1 and j_2 of module k is represented by two non-negative variables $\Delta_{k,j_1,j_2}^+, \Delta_{k,j_1,j_2}^- \geq 0$:

$$\Delta_{k,j_1,j_2}^+ - \Delta_{k,j_1,j_2}^- = \sum_{s=1}^S (b_{k,j_1,s} - b_{k,j_2,s}). \quad (6)$$

Thus, under above constraints we have to minimize

$$\sum_{k=1}^N \sum_{j_1=1}^{n_k-1} \sum_{j_2=j_1+1}^{n_k} (\Delta_{k,j_1,j_2}^+ + \Delta_{k,j_1,j_2}^- - \varepsilon_{k,j_1,j_2}). \quad (7)$$

Slack variables $0 \leq \varepsilon_{k,j_1,j_2} \leq \min\{D, \Delta_{k,j_1,j_2}^+ + \Delta_{k,j_1,j_2}^-\}$ allow absolute group size differences to vary between 0 and $D \in \{0, 1, 2, \dots\}$ without penalty.

Measures to speed-up the program (1)

Certain groups can be assigned to sub-slots in a fixed manner, for example if the number of sub-slots $m_k \cdot p_k$ equals the number of groups n_k . For such modules k we can set

$$a_{k,j,i,l} := \begin{cases} 1 & : j = (i - 1) \cdot p_k + l \\ 0 & : \text{otherwise.} \end{cases} \quad (8)$$

Measures to speed-up the program (2)

By assigning groups to sub-slots in a chronologically sorted manner due to their group number, one can avoid many permutations: For all modules $k \in [N]$, all time slots $i_1 \in [m_k]$ and all sub-slots T_{k,i_1,l_1} , $l_1 \in [p_k]$, and all groups $j_1 \in [n_k]$ let

$$\max \left\{ \sum_{i_2=i_1+1}^{m_k} \sum_{j_2=1}^{j_1-1} \sum_{l_2=1}^{p_k} a_{k,j_2,i_2,l_2}, \sum_{j_2=1}^{j_1-1} \sum_{l_2=l_1+1}^{p_k} a_{k,j_2,i_1,l_2} \right\} \\ \leq n_k \cdot (1 - a_{k,j_1,i_1,l_1}). \quad (9)$$

Practical application: summer semester 2019

Presented results belong to our bachelor programs in

- computer science (second and fourth semester)
 - 330 students including 59 dual education and part-time students,
 - 30 modules,
 - up to 8 groups per module
- electrical engineering (second, fourth, and sixth semester)
 - 168 students including 40 dual education and part-time students,
 - 27 modules,
 - up to 4 groups per module

Running times

slack size D	sorting (9)	init. (8)	computer science	electrical engineering
2	-	-	51.59	0.13
2	-	✓	3.35	0.1
2	✓	-	2.8	0.06
2	✓	✓	1.57	0.05
1	-	-	57.92	0.15
1	-	✓	6.32	0.08
1	✓	-	3.49	0.09
1	✓	✓	3.33	0.07
0	-	-	memory overflow	2.65

IBM ILOG CPLEX processor times (average mean of ten runs)
measured in seconds on an Intel Core i5-6500 CPU, 3.20 Ghz x4
with 16 GB RAM.

Enhancements

- The program can also be used as an online algorithm for adding late students.
- There might be no feasible solution if students enroll to conflicting modules of different semesters. Such situations can be identified before program execution.
- As a secondary optimization goal, one could maximize the number of students that get commonly assigned to groups along all modules.
- Students who have to repeat modules could be distributed as evenly as possible among groups, since experience has shown that for such students the risk of non-appearance is high.

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